VCS: A VIRTUAL COMPUTING AND STORAGE APPROACH TO COMPUTING CONTINUUM

Agenda

- Background and problem statement
- Existing solutions and limitation
- Topology and concept
- VCS overview
- Client architecture and awareness
- Virtual storage
- Computing continuum workflow
- Optimization of workflow
- Cloud-edge-client collaboration
- Test bed evaluation
- Future direction
- A very early demo

Background and problem statement

- Modern IT industry makes computer hardware cheaper
- 5G and broadband makes internet faster
- More users might have multiple computing devices like laptop, phone and special equipment
- New request comes out access same copy of desktop from multiple devices located at geographically dispersed locations

Existing solutions and limitation

	ZAP,	CR&TR,	JIT, Sledge	VDI	Application	Local-remote	User profile		
	NomadB	LS			cloudification	data mirroring			
What is migrated	process	VM	container	desktop	application	Data	profile	VM hosting @ cloud	App1 App2 Web hosting
	×	×	×	V	V	×	×		App1
Desktop continuum	×	V	×	V	×	×	V	Client A Client B	Browser A Br
No legacy	V	V	V	V	×	×	V	1) VDI	2) Application cloudifi
application porting								Remote Office365	Remote windows
No strong network	×	×	×	×	×	V	V	data	profile
dependency								Local data	windows
Computing	V	V	V	×	×	V	V	Office365 Office365	profile
performance near								OS OS	Windows V
to native								3) Local-remote data mirroring	4) User profile

- Steven et al proposed Zap, a pod or group of process migration solution based on an OS level thin virtualization layer.
- Jacob et al proposed an on-the-fly entire OS migration solution based on NomadBIOS, a Xen-based hypervisor with pre-copy migration to keep the OS running while migrating, tracking the changes of its address space and sending updates of image over a couple of iterations.
- Haikun et al provided a CR & TR (checkpointing/recovery and trace/replay) mechanism for fast and transparent VM migration design to reduce the migration time and save network bandwidth.
- Fei and Xiaoming et al proposed a three-layer image management mechanism and correlated layer moving system to improve the migration performance.
- Takahiro and Hirotaka et al suggested a very good live storage migration design with on-demand fetching and background copying strategy to keep IO virtual disk consistency while minimal IO performance impact.

Topology and concept





- Software stack, Desktop
- Computing tasks,
- Terminal and Client,
- Edge and cloud, Local,
- near and backend,
- Client awareness

VCS overview – problem statement and scope

• Problem statement:

- User could pick up any VCS terminal within network scope to access her private desktop environment
- VCS terminal is aware of software stack. It will designate the most suitable hardware nearby to provide the best user experience as native as PC.
- The software stack of each individual user is manageable, updatable, migratable and **propagatable**
- Scope:
 - VCS terminal is mostly PC-like, Different users will share similar software stack as much as possible
 - VCS works under **diversified network** environments
 - VCS would support all mainstream desktop OS
 - does not require instancy or live

Client architecture and awareness



1, bare-metal boot; 2, virtual machine boot; 3. boot with edge server

Awareness considerations:

- ISA/CPU arch
- Performance
- Commercial reasons



- Local bare-metal
- Local VM
- Adjacent edge

Virtual storage – image overlap problem



Result: conflict happens!

Virtual storage – personal computing engine



Virtual storage – an example

Private storage:

- Managed by each user ٠
- File and registry-based ٠ Only update when • write/update 0 0
- Files/registry in private ٠ has higher priority than User A that of system storage



0 0

Base +

delta#1





Delta #1

Base image

Sy	stem	stora	ge:	
•	N/0 m	agad	h.,	~

- administrator
- Block based storage ٠
- Incremental based ٠ 1,

Image	Operation
Base	Have 3 files: <mark>a, b, c</mark>
Delta #1	Update a to a', remove b. so base + delta#1 = <mark>a', c</mark>
Delta #2	Add file d, so base + delta # + delta#2 = <mark>a', c, d</mark>
User A private	Private file: a'', d''
User 2 private	Private file: <mark>e</mark>

	Combination	Result of files
	Base + User A	a'', b, c, d''
Base + delta#1 + delta#2	Base + delta #1 + user A	<mark>a'', c, d''</mark>
	Base + delta #1 + delta #2 + User A	<mark>a'', c, d''</mark>
	Base + User A	<mark>a, b, c, e</mark>
	Base + delta #1 + user A	a', c, e
	Base + delta #1+ delta #2 + User A	<mark>a', c, d, e</mark>

Conclusion: in any combination of system storage & private storage, private data (file) is still there no matter how system data is changed.

Computing continuum workflow



Optimization of workflow: the pre-copy mechanism



Cloud-edge-client collaboration



Test bed and evaluation



PassMark Rating for PC, VCS and VDI



Disk IO Benchmark - Single Thread





Future directions

- Application continuum
- Partial continuum?
- Non-network continuum
- Scalable cloud-edge-client hierarchy

A very early demo to show the VCS approach

THANK YOU